

Lecture 4

Cryptography



Mobile Business II (SS 2017)

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Introduction

- Symmetric Cryptosystems
- Public Key Cryptography

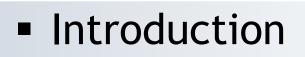


Cryptographic Systems

- Intention
 - Confidentiality (secrecy of messages): encryption systems
 - Integrity (protection from undetected manipulation) and accountability: authentication systems and digital signature systems
- Key distribution
 - Symmetric:

Both partners have the same key.

- Asymmetric: Different (but related) keys for encryption and decryption
- In practice mostly hybrid systems



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- Symmetric Cryptosystems
 - General Concept
 - Caesar Cipher
 - AES
 - Advantages and Problems
- Public Key Cryptography



Symmetric Encryption Systems

- Typical applications
 - confidential storage of user data
 - transfer of data between 2 users who negotiate a key via a secure channel
 - end-to-end channel encryption
- Examples
 - Vernam-Code (one-time pad, Gilbert Vernam)
 - key length = length of the plaintext (information theoretically secure)
 - DES: Data Encryption Standard
 - key length 56 bit \rightarrow 2⁵⁶ different keys
 - AES: Advanced Encryption Standard (Rijndael, [NIST])
 - 3 alternatives for key lengths: 128, 192 and 256 bit

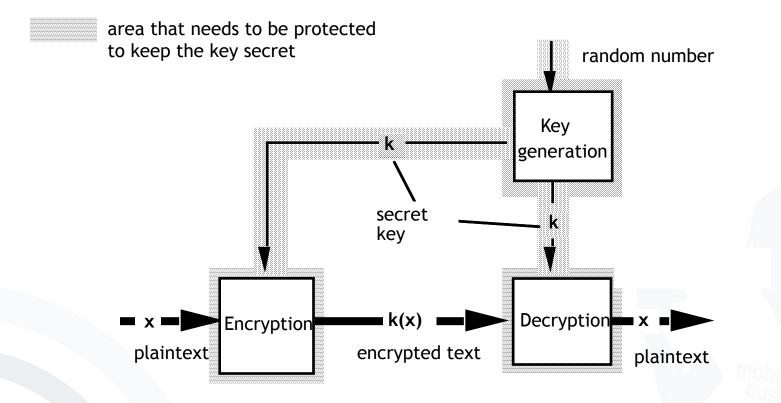
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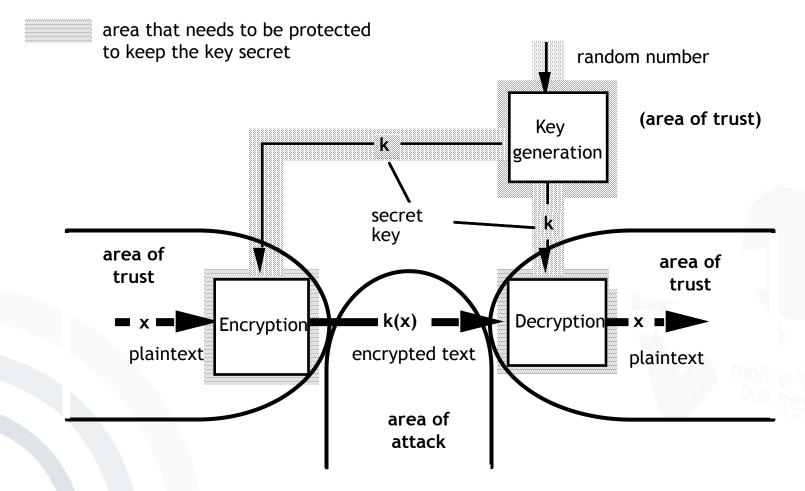
Symmetric Encryption Systems (1)



black box with lock, two equal keys

[based on Federrath and Pfitzmann 1997]

Symmetric Encryption Systems (2)



[based on Federrath and Pfitzmann 1997]



)

- Keys have to be kept secret (secret key crypto system).
- It must not be possible to derive the plaintext or the used keys from the encrypted text (ideally encrypted text is not distinguishable from a numerical random sequence).
- Each key shall be equally probable.
- In principle each system with limited key length is breakable by testing all possible keys.
- Publication of encoding and decoding functions (algorithms) is considered as good style and is trust-building.
- Security of cryptosystems should base on the strength of chosen key lengths.

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Caesar Cipher



18 19 20 21 22 23 24 25

We assign a number for every character.

17

16

15

14

 This enables us to calculate with letters as if they were numbers.

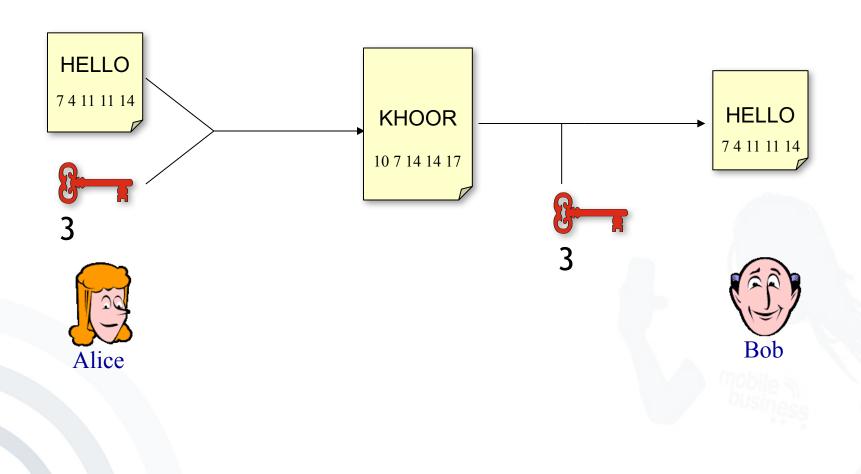




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Caesar Cipher: Example







- Very simple form of encryption.
- The encryption and decryption algorithms are very easy and fast to compute.
- It uses a very limited key space (n=26)
- Therefore, the encryption is very easy and fast to compromise.

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- The Data Encryption Standard (DES) was designed to encipher sensitive but not classified data.
- The standard has been issued in 1977.
- In 1998, a design for a computer system and software that could break any DES-enciphered message within a few days was published.
- By 1999, it was clear that the DES no longer provided the same level of security it had 10 years earlier, and the search was on for a new, stronger cipher.
- The successor is called Advanced Encryption Standard (AES).
- AES has been approved for Secret or even Top Secret information by the NSA. [Bishop 2005]

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Symmetric Encryption



Advantage: Algorithms are very fast

Algorithm	Performance*
RC6	78 ms
SERPENT	95 ms
IDEA	170 ms
MARS	80 ms
TWOFISH	100 ms
DES-ede	250 ms
RIJNDEAL (AES)	65 ms

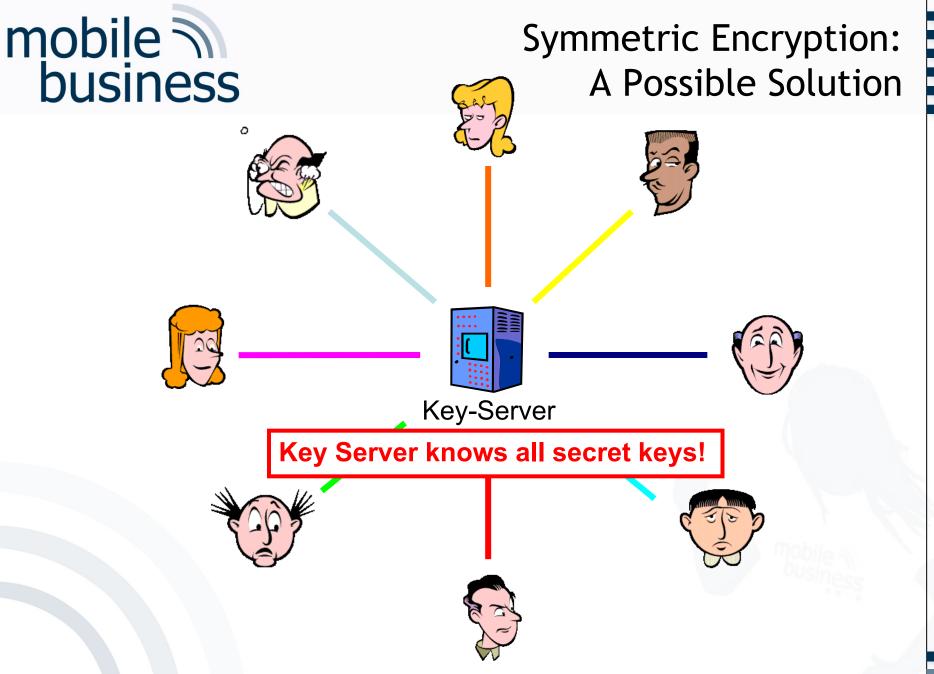
* Encryption of 1 MB on a Pentium 2.8 GHz, using the FlexiProvider Java)

[J. Buchmann: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]

Problems of Symmetric Cryptosystems: Key Exchange

n*(n-1)/2 Keys Internet: ~ 1.000.000.000 Users => ~ 500.000.000.000.000 Keys [J. Buchmann: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]

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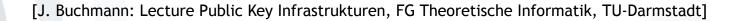


[J. Buchmann: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]

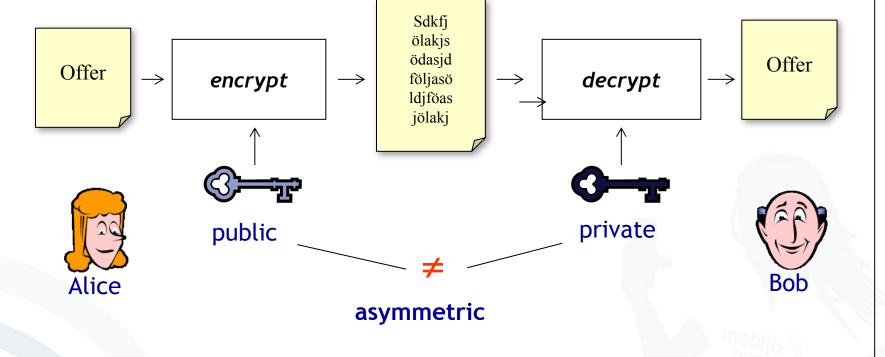


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- Public Key Cryptography
 - General Concept
 - Algorithms
 - Hybrid Systems
 - Digital Signature
 - Key Management
 - Example: PGP

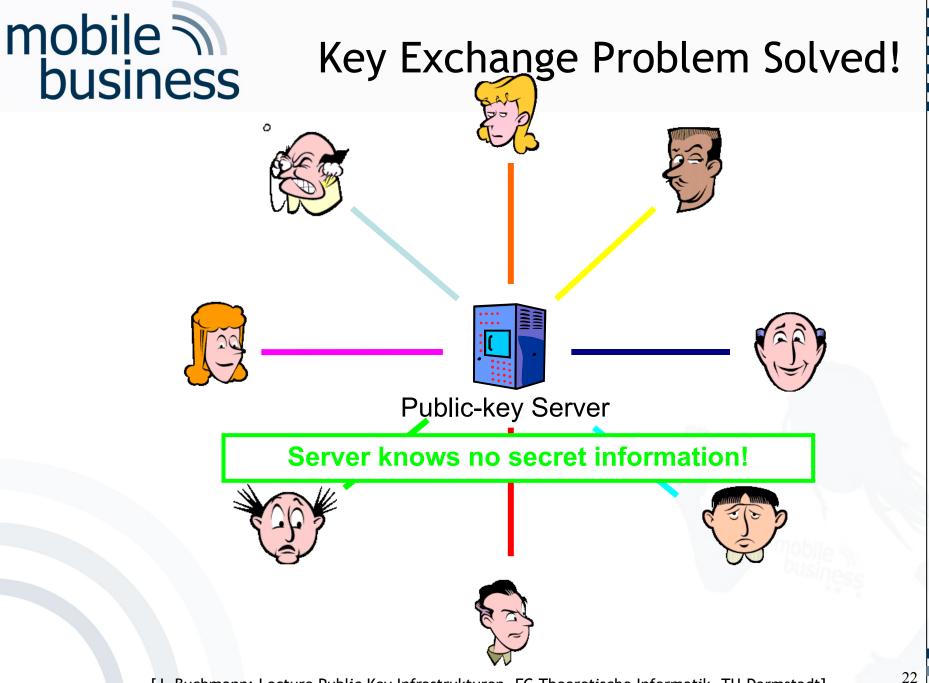




Public Key Encryption







[J. Buchmann: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]



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Concept of Asymmetric Encryption Systems

- Use of 'corresponding' key pairs instead of one key:
 - **Public key** is **solely** for encryption.
 - Encrypted text can only be decrypted with the corresponding private (undisclosed) key.
- Deriving the private key from the public key is hard (practically impossible).
- The public key can be distributed freely, even via insecure ways (e.g. directory (public key crypto system)).
- Messages are encrypted via the public key of the addressee.
- Only the addressee possesses the private key for decoding (and has to manage the relation between the private and the public key).

mobile business **Asymmetric Encryption Systems** random number area that needs to be protected to keep the key secret Key geneencryption key, ration publicly known decryption key,

plaintext Encryption Encrypt

box with slot, access to messages only with a key

[based on Federrath and Pfitzmann 1997]



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Asymmetric Encryption Systems: Examples



RSA

- Rivest, Shamir, Adleman, 1978
- is based on the assumption that the factorization of the product of two (big) prime numbers (p*q) is "difficult" (product is basis for the keys)
- key lengths typically 1024 bit, today rather 2048

[Rivest et al., 1978]

Diffie-Hellman

- Diffie, Hellman, 1976, first patented algorithm with public keys
- allows the exchange of a secret key
- is based on the "difficulty" of calculating discrete logarithms in a finite field

[Diffie, Hellman, 1976]





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Performance of Public Key Algorithms

Algorithm	Performance*	Performance compared to Symmetric encryption (AES)			
RSA (1024 bits)	6.6 s	Factor 100 slower			
RSA (2048 bits)	11.8 s	Factor 180 slower			

Disadvantage: Complex operations with very big numbers

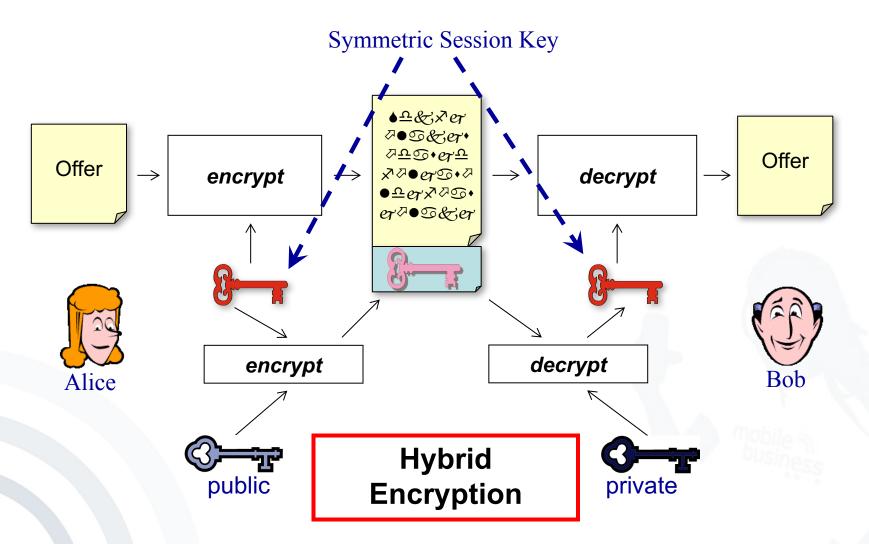
\Rightarrow Algorithms are very slow.

* Encryption of 1 MB on a Pentium 2.8 GHz, using the FlexiProvider (Java)

[J. Buchmann: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]



Solution: Hybrid Systems



[based on: J. Buchmann 2005: Lecture Public Key Infrastrukturen, FG Theoretische Informatik, TU-Darmstadt]



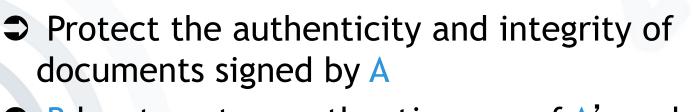
Introduction

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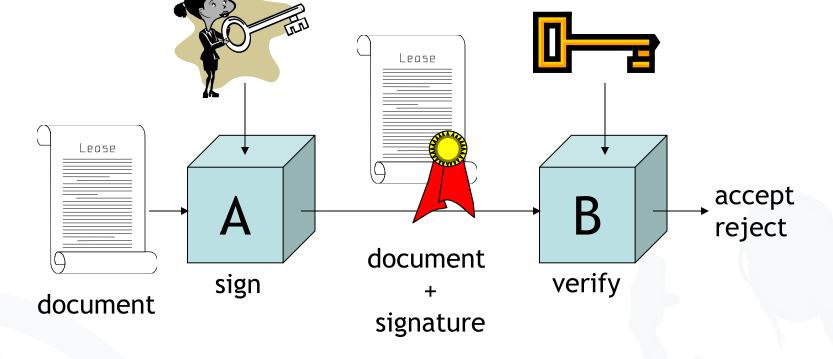
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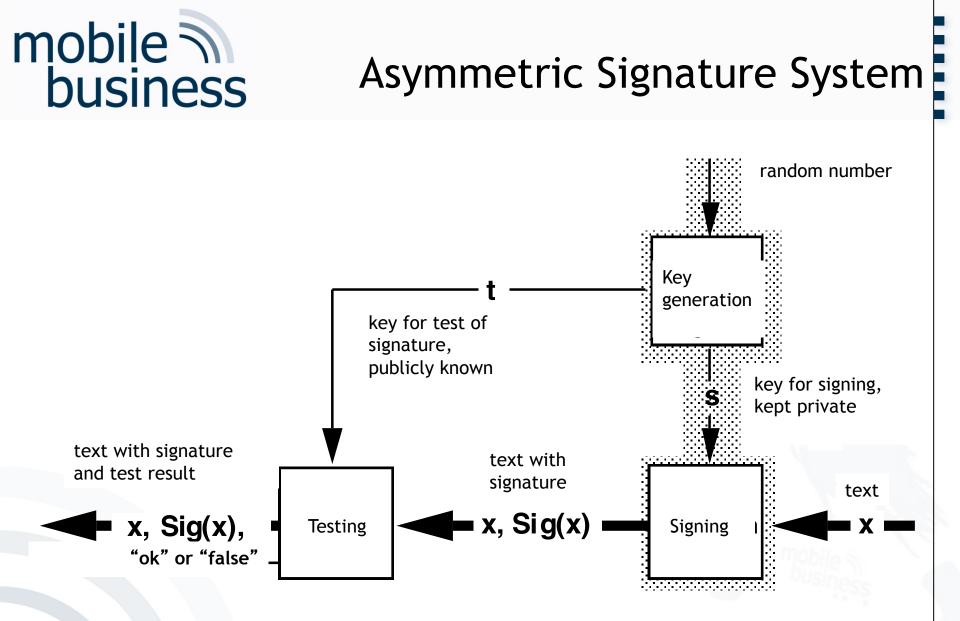
Digital Signatures



 \bigcirc B has to get an authentic copy of A' s public key.







Iocked glass show-case; just one key to put something in

[Federrath and Pfitzmann 1997]

Example PGP: Encrypt and Sign a Message

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Example PGP: Decrypt and Check a Message

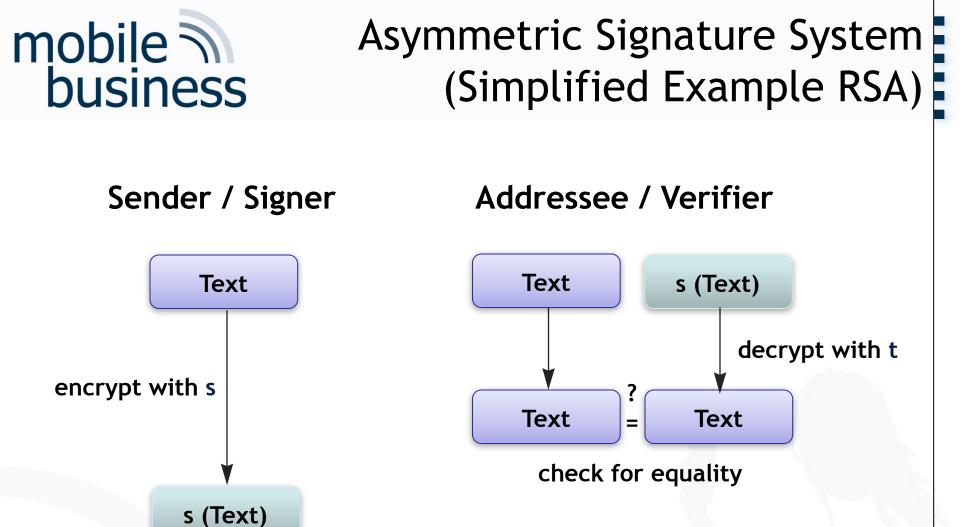
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			Heiko Rossnagel <heiko.rossnagel@m-lehrstuhl.de> (DH/2048)</heiko.rossnagel@m-lehrstuhl.de>			
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*** END PGP DECRY	*** END PGP DECRYPTED/VERIFIED MESSAGE ***					
			*			
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Asymmetric Signature Systems: Examples

- RSA: Rivest, Shamir, Adleman
 - Asymmetric encryption system which also can be used as a signature system via "inverted use",
 - Message encrypted with the private key (= signing key) gives the signature,
 - Decoding with the public key (=testing key) has to produce the message.

[Rivest et al. 1978]

- DSA: Digital Signature Algorithm
 - Determined in the Digital Signature Standard of the NIST (USA),
 - Based on discrete logarithms (Schnorr, ElGamal),
 - Key length is set to 1024 bit.



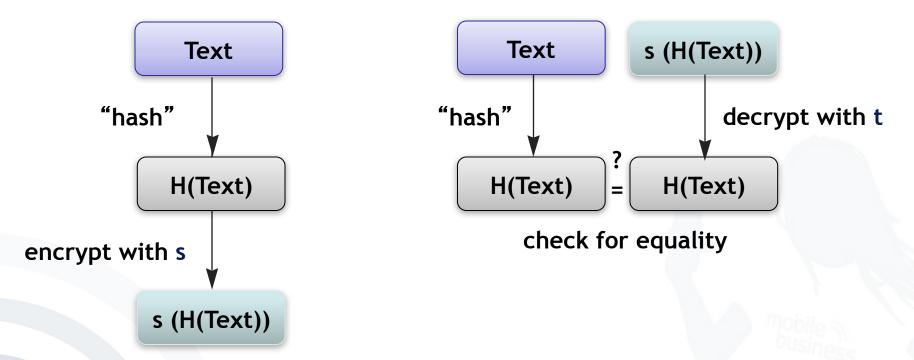
Signing key s only with the sender, test key t public
 Example is often mistakenly generalized.



Asymmetric Signature System (Example RSA)

Sender / Signer





Signing key s only with the sender, test key t public
 Example is often mistakenly generalized.





Hash Functions

- General hash functions (H(s))
 - Transformation of an input string s into an output string h of fixed length which is called hash value.
 - Example: mod 10 in the decimal system
- Cryptographic hash functions
 - Generally require further characteristics
 - H(s) is easily to compute for each s.
 - *H*(*s*) must be difficult to invert: In terms of figures it is difficult to compute *s* from *h*.
 - Virtual collision freedom: In terms of figures it is difficult to create collisions H(s1) = H(s2).
 - Examples: SHA-1, MD5, MD4



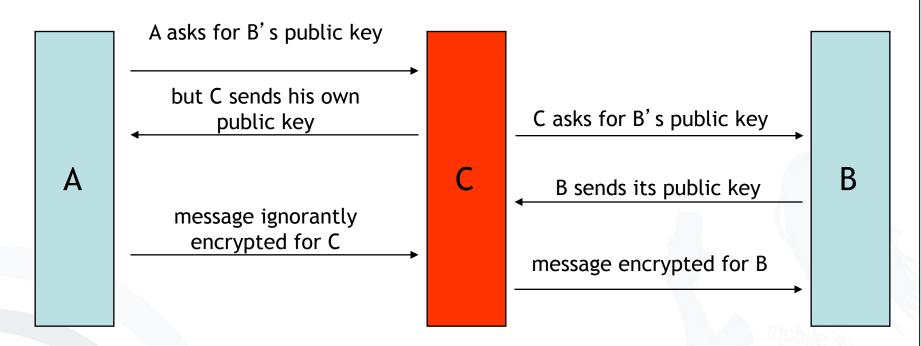
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"Man in the middle attack"



Keys are certified: a 3rd person/institution confirms (with its digital signature) the affiliation of the public key to a person.

Certification of Public Keys (1)

- **B** can freely distribute his own public key.
- But: Everybody (e.g. C) could distribute a public key and claim that this one belongs to B.
- If A uses this key to send a message to B, C will be able to read this message!

Thus:

How can **A** decide if a public key was really created and distributed by **B** without asking **B** directly?

- Keys get certified, i.e. a third person/institution confirms with its (digital) signature the affiliation of a public key to entity B.
- Public Key Infrastructures (PKIs)

Certification of Public Keys (2)

Three types of organization for certification systems (PKIs?):

- Central certification authority (CA)
 - A single CA, keys often integrated in checking software
 - Example: older versions of Netscape (CA = Verisign)
- Hierarchical certification system
 - CAs which in turn are certified by "higher" CA
 - Examples: PEM, Teletrust, infrastructure according to Signature Law

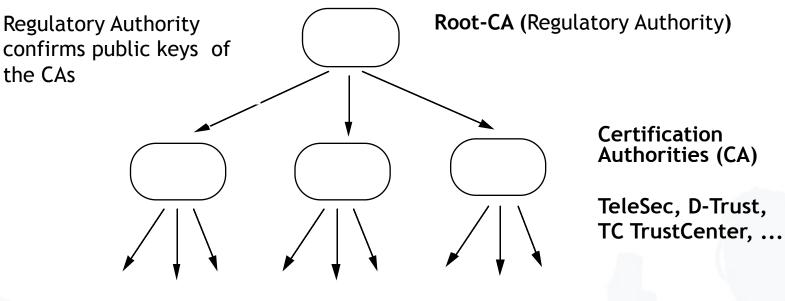
Web of Trust

- Each owner of a key may serve as a CA
- Users have to assess certificates on their own
- Example: PGP (but with hierarchical overlay system)



Hierarchical Certification of Public Keys

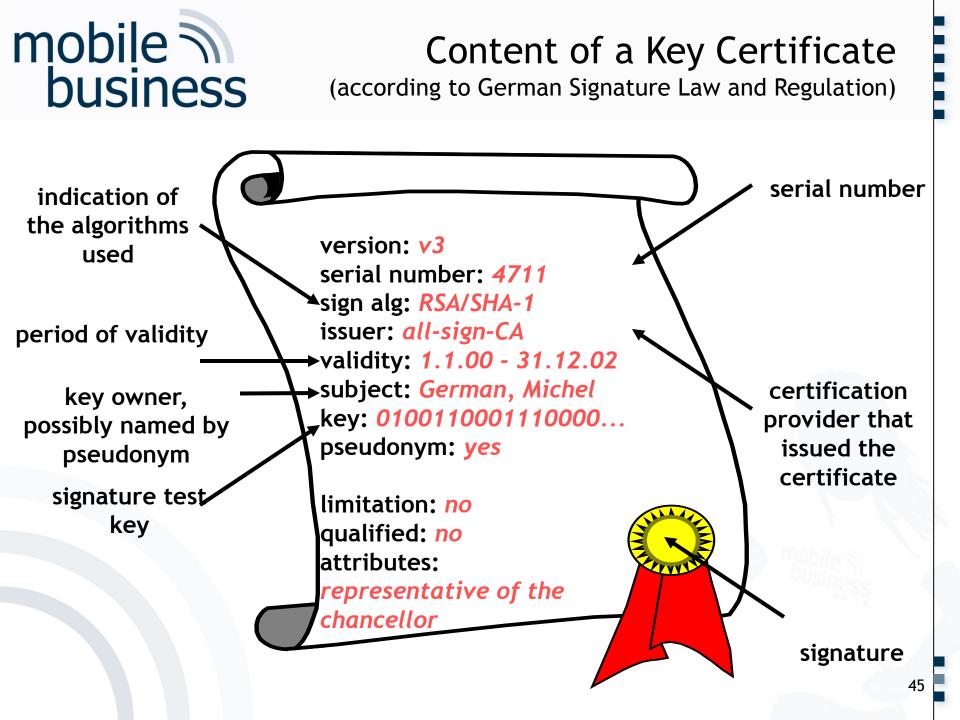
(Example: German Signature Law)



persons

organizations

- The actual checking of the identity of the key owner takes place at so called Registration Authorities (e.g. notaries, bank branches, T-Points, ...)
- Security of the infrastructure depends on the reliability of the CAs.







- Reliable identification of persons who apply for a certificate
- Information on necessary methods for fraud resistant creation of a signature
- Provision for secure storage of the private key
 - at least Smartcard (protected by PIN)
- Publication of the certificate (if wanted)
- Barring of certificates
- If necessary issuing of time stamps
 - for a fraud resistant proof that an electronic document has been at hand at a specific time





- Checking of the following items by certain confirmation centers (BSI, TÜVIT, ...)
 - Concept of operational security
 - Reliability of the executives and of the employees as well as of their know-how
 - Financial power for continuous operation
 - Exclusive usage of licensed technical components according to SigG and SigV
 - Security requirements as to operating premises and their access controls

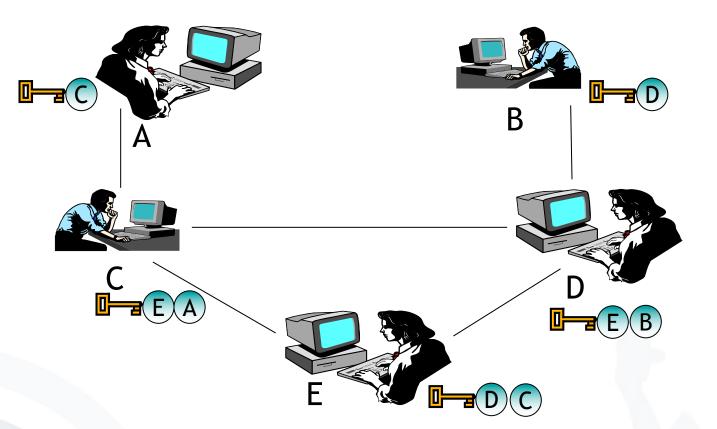
Possibly license of the regulation authority

mobile business Web of Trust "Introducer" David Bob knows David and has received David's public key by David himself Alice lets David sign her public key Bob can verify Alice' key Alice sends the signed on the basis of David's key to Bob signature Bob encrypts his message to Alice Bob Alice with the received key 5

- Each user can act as a "CA".
- Mapping of the social process of creation of trust.
- Keys are "certified" through several signatures.
- Expansion is possible by public key servers and (hierarchical) CAs.



Web of Trust Example



Web of Trust:

- Certification of the public keys mutually by users
- Level of the mutual trust is adjustable.



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Protection of Email Example PGP

- PGP = Pretty Good Privacy
 - De facto-Standard for freely accessible e-mail encryption systems on the Internet
 - First implementation by Phil Zimmermann
 - Long trial against Phil Zimmermann because of suspicion of violation of export clauses
 - In U.S., free version in cooperation with MIT (agreement with RSA because of the patent)
 - Meanwhile commercialized: www.pgp.com
 - Gnu Privacy Guard (GPG): non-commercial Open Source variant (OpenPGP, RFC2440)

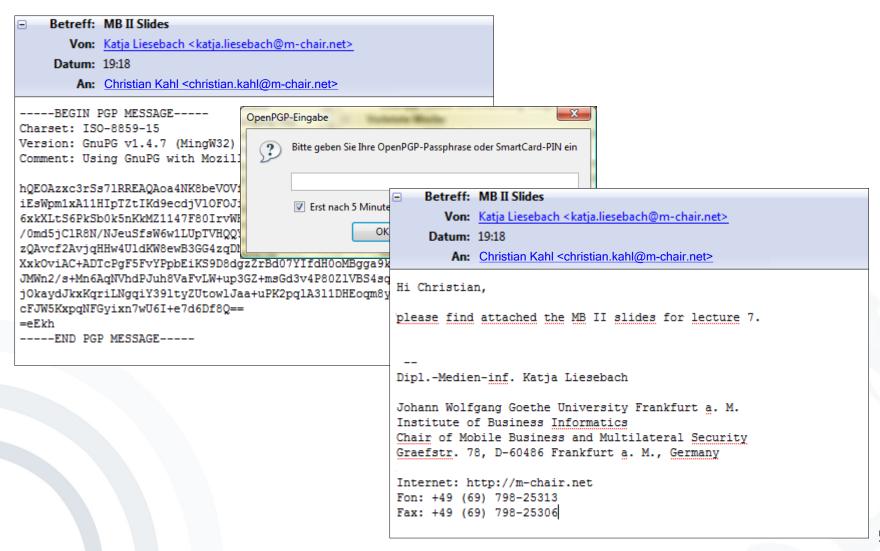
PGP: Encrypt Message

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	Conventional Encryption			

OpenPGP: Encrypt Message

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, _				-			
∃ Christian Kahl < christian.kahl@m-le	Nicht	gefundene Empfänger					
	Empfä	änger für Verschlüsselung wählen					
	Be	Benutzer-ID		Vertrauen	Ablauf	Schlüssel-ID	EŞ.
		Christian Kahl < christian.kahl@m-lehrstuhl.de>		absolutes Ver		14E21EDA	
		Alexander Boettcher ("Nur wenige wissen, wie viel man wis	sen muss, um zu	abgelaufen	02.09.2006	8D539C6E	·····
ched the MB II slides for lect		Alexander Boettcher <ab764283@inf.tu-dresden.de></ab764283@inf.tu-dresden.de>		-		A63325B3	
		Alexander Boettcher <ab764283@os.inf.tu-dresden.de></ab764283@os.inf.tu-dresden.de>		abgelaufen	11.10.2005	F26EE0CD	
		Andre Meixner <s4538672@inf.tu-dresden.de></s4538672@inf.tu-dresden.de>		-		7C433232	
OpenPGD Restätigung		× · · · · · · · · · · · · · · · · · · ·		-		7E39E652	-
OpenPOP-bestatigung			ELEVANCE	-		52B1B05D	-
VERSCHLÜSSELTE Nachricht an	n folgende F	Empfänger senden:		-			
50		roigende Emphanger senden.					
christian.kahl@m-lehrstuhl.de				-			
				-	11.04.0014		
Hinweis: Die Nachricht wurde n	nit folgende	len Benutzer-IDs / Schlüsseln verschlüsselt:	e>	-	11.04.2011		
				-	20.02.2000		
	- L-			-	20.02.2009		
	Īa			-			
		Natia Liesepach < katia.iiesepach@m-chair.net>		absolutes Ver			
				-		F7C207CE	-
) Katrin Borcea <kati@inf.tu-dresden.de></kati@inf.tu-dresden.de>		-			
	Nac	chricht unverschlüsselt und nicht unterschrieben senden					
	Dies	esen Dialog nicht mehr anzeigen, wenn Verschlüsselung unr	nöglich ist				
	Liste ak	ktualisieren Fehlende Schlüssel herunterladen					
	Liste di					_	
					OK	Abbrec	hen
	Rechtschr. Anhang OpenP Akatja.liesebach@m-chair.net> - kation Christian Kahl < christian.kahl@m-le	Rechtschr. Anhang OpenPGP <katja.liesebach@m-chair.net> - katjobach@m-Chair.net> Christian Kahl < christian.kahl@m-le</katja.liesebach@m-chair.net>	Rechtschr. Anhang OpenPGP S/MIME Speichern c katja.liesebach@m-chair.net> - katjette te	Rechtschr. Anhang OpenPGP S/MIME Speichern Isektiga.liesebach@m-chair.net> kati-liesebach@m-chair.net> OpenPGP-Schlüssel auswählen Christian Kahl < christian.kahl@m-le	Rechtschr. Anhang OpenPGP S/MIME Speichern a <katja.liesebach@m-chair.net> - katja.liesebach@m-chair.net> - katja.liesebach@m-chair.net> - katja.liesebach@m-chair.net> Christian Kahl <christian.kahl@m-let< td=""> PenPGP-Schlüssel auswählen - Nicht gefundene Empfänger Empfänger für Verschlüsselung wählen - Wertrauen - Wertrauen Wertrauen - Wertrauen - Wertrauen Wertrauen</christian.kahl@m-let<></katja.liesebach@m-chair.net>	Rechtschr. Anhang OpenPGP SVMME Speichern • katjaliesebach@m-chair.net> - katjaliesebach@m-chair.net> - katjaliesebach@m-chair.net> - katjaliesebach@m-chair.net> • Christian Kahl < christian.kahl@m-leh	Rechtschr. Anhang OpenPGP S/MME Speichern • katjaliesebach@m-chair.ne>

OpenPGP: Decrypt Message



PGP: Decrypt Message

etreff: Klausur MC1	Cc		-
BEGIN PGP MESSAGE		-	
Version: PGP 8.0 - not lic	ensed for commercial use: www.pgp.co	<u>m</u>	
aHEsO7/tFrJFQJpPBcUWouy47p B5TXKtUB8YJdpPnck61as78RBP Ag4DIY1owhVX6ZwQCAD2L9WAA9 Alkh23iQO1I9Drye/uygpcQpT2 cDP3GEanyDiDU6R9F1XFOvxPNM aOuXNA9iAC96dhg7NpvzCJI2J7 dfvQ3NiGrUEQsOHVxwjQdMtr8C i77MitBfAbxXF0gFS7/b2Lccba h2oTOSjWCRp/v5s9Og1aUtcAxd n39jRjPE9Ob/HLjMwPAXUHyneh	Sk4G/TAexpMLX436biwBp6xP8pa89R7ro5Xo 4sR2FO+IXqJuJyHp5ExMGIdmQCpGXEoS2Ijw 1sq8VDrAlYopEAeqMMw2pkBuoxyo3KCiRkhi 7xEUBWMET6kR9n5+oafTBF+R0lv6UOz2TO55 HhTtZYlAjjudLvi+GsegO1WmBjY8q8G1Y61C k6Ek8hH6qZ37hhDNDCXkxkSjM3nJ2VuuLvXb xRMtuBc9BUI8LX0DrvGLwnLtaD5+EvgL1xTu 09kREYLuAdD7j/05WtsAdbAVMn72PYF0IRfZ K8fx6e1VNFnVO7B/9qpd0Gg5WZVP2eQA5fbw 1RAjQPHpVsFS2eXXMnC9Z2vNIFMh6Ktqmt6E 9QrCX1XSqH0RNcjIYVrnQyZGIk8t39059FBd	Heiko Rossnagel <heiko.rossnagel@m-lehrstuhl.de> (Jan Muntermann <munterma@wiwi.uni-frankfurt.de> (Enter passphrase for your private key :</munterma@wiwi.uni-frankfurt.de></heiko.rossnagel@m-lehrstuhl.de>	DH/2048)
	EVRC6piJaJFmrzifnzliwfuf82Tc42GBd9bP x4nGFX17Hy1vjkqm7gfyLxXjgeDCnjxm7O8J		
JjwtR+1SkqMCXs+PzcAHDsiuG:		OK	Cancel
E3huhK5cfvu1Ug7+Oa9SUAy4	Text Viewer		
NZncI3vJgkZeZrlbh+pi4dRjs(=hCO9 END PGP MESSAGE	Hello Jan		
heiko rossnagel frankfurt direkt -25306 D-60054 frankfurt	Please find the MB 2 exam enclosed.		lie S
	Copy to Clipboa	rd OK	

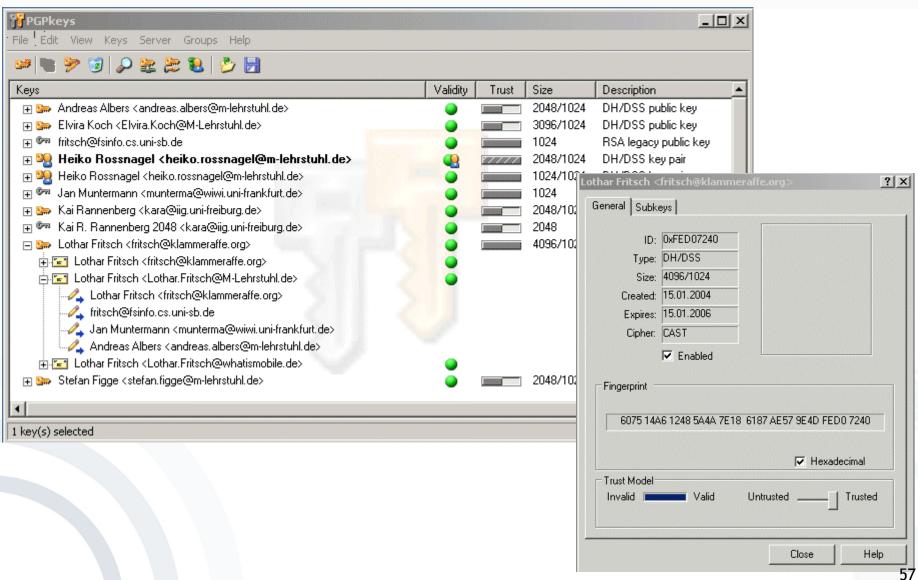


PGP-Certification of Keys

- Certification of public keys by users: "Web of Trust"
- Differentiation between 'validity' and 'trust'
 - 'Trust': trust that a person / an institution signs keys only if their authenticity has really been checked
 - 'Validity': A key is valid for me if it has been signed by a person / an institution I trust (ideally by myself).
- Support through key-servers:
 - Collection of keys
 - Allocation of 'validity' and 'trust' remains task of the users
- Path server:

Finding certification paths between keys

PGP: Key Management



OpenPGP: Key Management

eige Schlüssel, deren Benutzer-ID oder Schlüssel-ID folgendes enthalten:			Alle <u>z</u> eig	en			
Benutzer-ID Alexander Boettcher ("Nur wenige wissen, wie viel man wissen muss, um z Alexander Boettcher <ab764283@inf.tu-dresden.de> Alexander Boettcher <ab764283@os.inf.tu-dresden.de> Andre Meixner <s4538672@inf.tu-dresden.de></s4538672@inf.tu-dresden.de></ab764283@os.inf.tu-dresden.de></ab764283@inf.tu-dresden.de>	Vertrauen abgelaufen absolutes Vertr abgelaufen -	11.10.2005	öffentli öffentli öffentli öffentli	ch ch			
Andreas Albers <andreas.albers@m-lehrstuhl.de> Andreas Pfitzmann <pfitza@inf.tu-dresden.de> NO LEGAL RELEVANCE André Deuker <andre.deuker@m-lehrstuhl.de> Birgit Pretscheck <birgit.pretscheck@gmx.net></birgit.pretscheck@gmx.net></andre.deuker@m-lehrstuhl.de></pfitza@inf.tu-dresden.de></andreas.albers@m-lehrstuhl.de>	absolutes Vertr absolutes Vertr absolutes Ve		öffentli öffentli haften			(Internet)	
Christian Kahl < christian.kahl@m-lehrstuhl.de> Denis Royer < me@myasterisk.de> Elvira Koch < Elvira.Koch@m-lehrstuhl.de> Felix Göpfert (keine Passphrase) < fg798936@inf.tu-dresden.de> Hagen Wahrig < wahrig@web.de> Jan Zibuschka < zibuschka@m-lehrstuhl.de> Kati Rannenberg < Kai.Rannenberg@m-lehrstuhl.de> Katja Liesebach < katja.liesebach@inf.tu-dresden.de> Katja Liesebach < katja.liesebach@m-chair.net> Katrin Borcea < kati@inf.tu-dresden.de> Marco Lehmann < m99@gmx.li>	absolutes Ve absolutes Ve - absolutes Ve absolutes Ve - absolutes V e -	Primäre Benut Schlüssel-ID Typ Vertrauen Besitzer-Vertra Fingerabdruck	[[] uen	Christian Kahl < 0x14E21EDA öffentlich absolutes Vertra absolutes Vertra E1CC 3AA5 BCE	uen		rstuhl.de> B8 B299 14E2 1EDA
Marco Lemmann (mb)@gma.n> Mathias Staab < mathias.staab@arcor.de> Mike Bergmann (dienstlich. TU Dresden. unbeschrnkt αltiα) < mb41@inf.t		Typ Unterschlüs	ID sel 0x5	Algo 98F0 ELG	Stär 2048	Erzeugt 07.09.2007	Ablauf-Datum nie



Key-Server

🎢 PGPkeys Search Window					-DX
Search for keys on Idap://keyserver.pgp.com 💌 where					Search
User ID 💽 contains 💌 Kai Rannenberg					Clear Search
More Choices Fewer Choices Search Pending Area					Help
Keys	Validity	Trust	Size	Description	
🕀 🖙 Kai R. Rannenberg 2048 <kara@iig.uni-freiburg.de></kara@iig.uni-freiburg.de>	0		2048	RSA legacy	public key
🕀 🖙 Kai R. Rannenberg <kara@iig.uni-freiburg.de></kara@iig.uni-freiburg.de>			1024	RSA legacy	public key
🕀 🖙 kara <kara@iig.uni-freiburg.de></kara@iig.uni-freiburg.de>	0		2048/1024	DH/DSS put	olic key
🖆 Found 3 key(s) matching search criteria.					1.

Sch

OpenDGD Schlücsel herunterladen

	Ope	nPOP-Schlussel heruntenaden			_
	ſ	Schlüssel gefunden - Auswählen zum Importieren			
		Benutzer-ID	Erstellt	Schlüssel-ID	C‡
		□	1997-03-09	8EF041F1	
		□	1997-09-18	AF1FDF70	
		□	1997-03-09	6623E0DD	
lüssel-Server auswählen					
uche nach Schlüssel Kai	Rannenberg				
chlüssel-Server subl	keys.pgp.ne	t v			
	5 1 5				
ſ	ОК	Abbrechen			
			0	K Abbre	echen



57

PGP: Public Key Catalogs

 File Edit		ose Index ``Kai Rannenberg '' - Microsoft Internet Explorer 📃 🗖
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Address 🥻	http://blackhole.pc	a.dfn.de:11371/pks/lookup?op=vindex&search=Kai+Rannenberg 🗸 🖉 Go 🛛 Link
.= <u>1</u> =		
Puł	blic Key	' Server Verbose Index ``Kai Rannenberg ''
Type b)its/keyID	Date User ID
		1997/09/18 kara <kara@iig.uni-freiburg.de></kara@iig.uni-freiburg.de>
sig	0B6375FD	Matthias Schunter <schunter@acm.org></schunter@acm.org>
sig	D5CDE083	Herbert Damker <damker@iig.uni-freiburg.de></damker@iig.uni-freiburg.de>
sig	879AC041	Birgit Pfitzmann 1 < <u>pfitzb@informatik.uni-hildesheim.de</u> > NO LEGAL RELEVANCE
sig	8128DC75	Gerhard Weck <73064.2271@compuserve.com>
sig	8EF041F1	Kai R. Rannenberg 2048 < <u>kara@iig.uni-freiburg.de</u> >
sig	2F8D5039	Kai Martius <kai@imib.med.tu-dresden.de></kai@imib.med.tu-dresden.de>
sig	5C3C4FE4	Holger Reif < <u>reif@prakinf.tu-ilmenau.de</u> >
sig	AF1FDF70	kara <kara@iig.uni-freiburg.de></kara@iig.uni-freiburg.de>
sig	49EF1D84	Hannes Federrath < <u>federrath@inf.tu-dresden.de</u> >
		Kai R. Rannenberg < <u>kair@microsoft.com</u> >
sig	<u>OB6375FD</u>	Matthias Schunter < <u>schunter@acm.org</u> >
sig	AEB4BCDD	fapp2_AEB4BCDD_HSK < <u>fapp2@cam.ac.uk</u> >
sig	AF1FDF70	kara < <u>kara@iig.uni-freiburg.de</u> >
sig	<u>044584B5</u>	Douglas Swiggum < <u>Swiggum@Waisman.Wisc.Edu</u> >
		Kai Rannenberg < <u>kara@iig.uni-freiburg.de</u> >
sig	OCB6E63F	Martin Reichenbach < <u>marei@iig.uni-freiburg.de</u> >
sig	AF1FDF70	kara < <u>kara@iig.uni-freiburg.de</u> >
		kara < <u>kara@telematik.iig.uni-freiburg.de</u> >
sig	<u>OB6375FD</u>	Matthias Schunter < <u>schunter@acm.org</u> >
sig	AF1FDF70	kara < <u>kara@iig.uni-freiburg.de</u> >
		Kai R. Rannenberg < <u>kara@iig.uni-freiburg.de</u> >
sig	<u>0B6375FD</u>	Matthias Schunter < <u>schunter@acm.org</u> >
sig	AF1FDF70	kara < <u>kara@iig.uni-freiburg.de</u> >
•		
e)		internet

- Network of public-key servers:
 - www.cam.ac.uk.pgp.net/pgpnet/email-key-server-info.html
 - http://pgp.mit.edu/





PGP: Practical Attacks and Weaknesses

- Brute-Force-Attacks on the pass phrase
 - PGPCrack for conventionally encrypted files
- Trojan horses, changed PGP-Code
 - e.g. predictable random numbers, encryption with an additional key
- Attacks on the computer of the user
 - Not physically deleted files
 - Paged memory
 - Keyboard monitoring
- Analysis of electromagnetic radiation
- Non-technical attacks
- Confusion of users [Whitten, Tygar 1999]

Remark

"Anybody who asserts that a problem is readily solved by encryption, understands neither encryption nor the problem."

(Roger Needham / Butler Lampson)

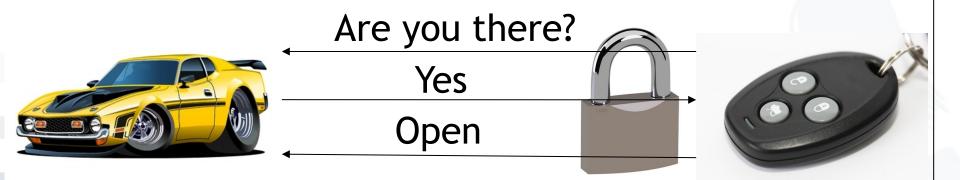
[Marshall Symposium 1998] [Randell 2004]







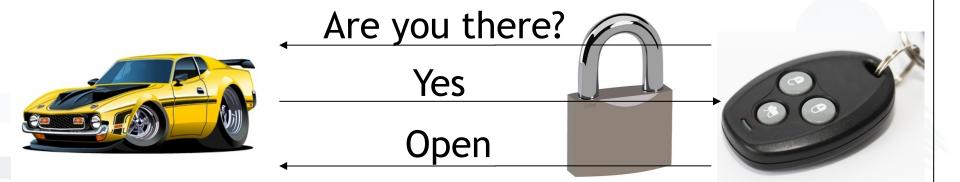
- Solution: Protect communication with crypto?
- e.g. symmetric cryptography + hash/signature

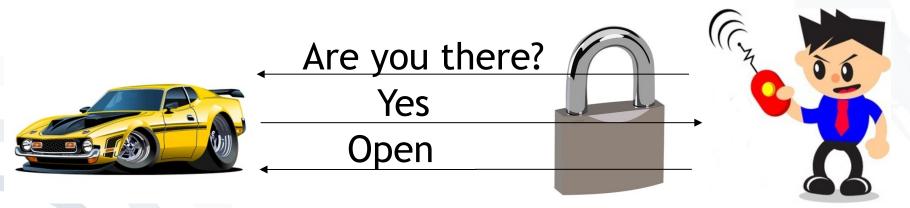


Replay Attack: Eavesdrop

mobile business



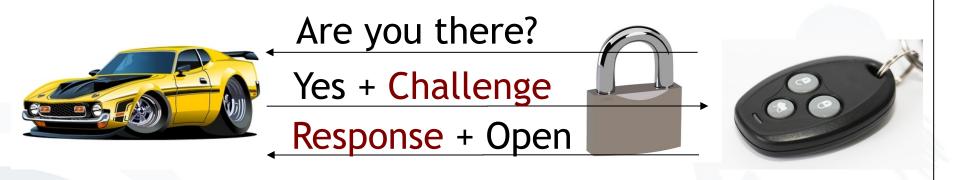


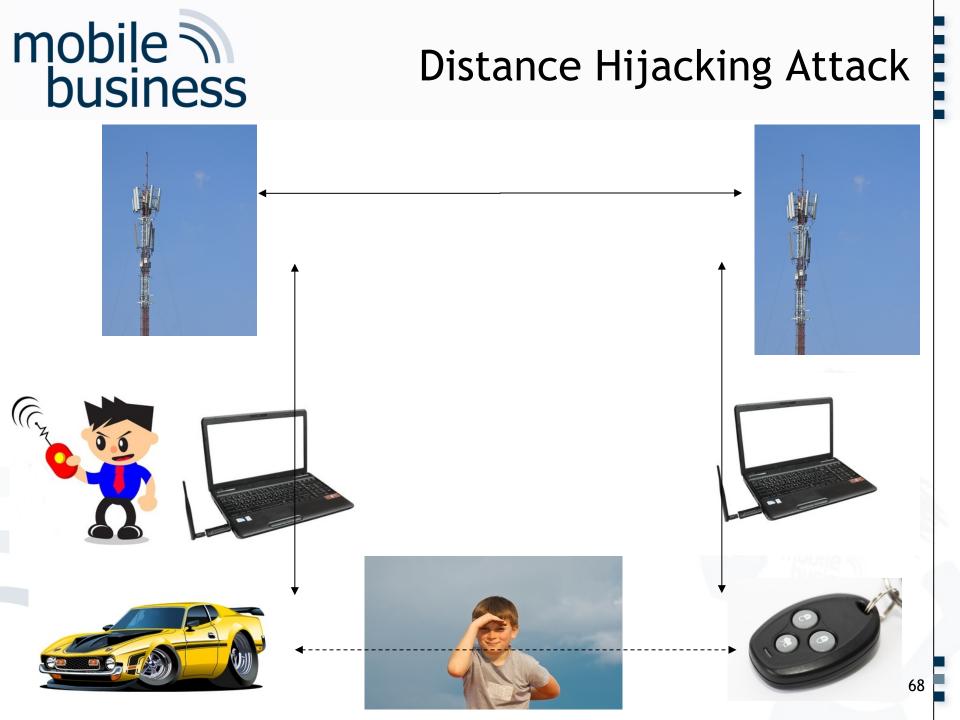




Replay Attack: Solution

e.g. Challenge-Response helps







Distance Hijacking Attack: Solution

Distance-Bounding Protocol









Side-Channel Attacks I

 A secure cryptoalgorithm does not imply that the implementation is also secure

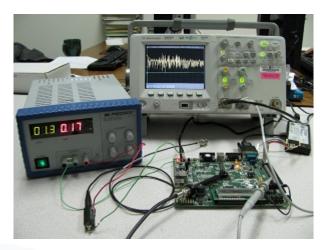


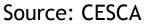
Source: Eran Tromer



Side-Channel Attacks II

Side-Channels: Time, Power, Noise, Radiation, ...





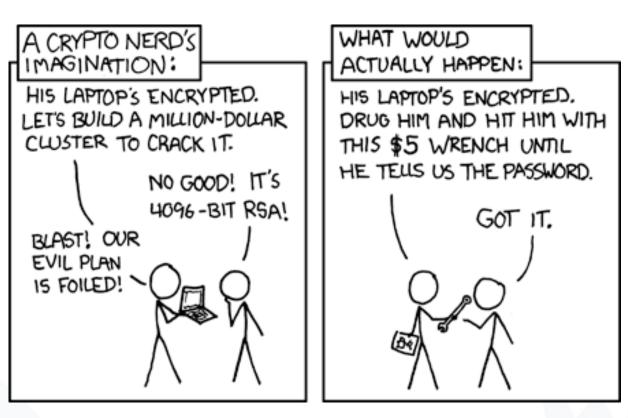


Source: Gilbert Goodwill

- Other data (side-channel) leaks information
- Conclusion on processed bits possible



The Human Element



Source: https://xkcd.com/538/

Human Element: Behavior and Passwords

- Florencio, D. & Herley, C., 2007. A large-scale study of web password habits. *Proceedings of the 16th international conference on World Wide Web - WWW '07*, p.657. Available at: http://portal.acm.org/citation.cfm?doid=1242572.12426 61.
- Florêncio, D., Herley, C. & Coskun, B., 2007. Do strong web passwords accomplish anything? *Proceedings of the* 2nd USENIX workshop on Hot topics in security (HOTSEC'07), p.10. Available at: http://portal.acm.org/citation.cfm?id=1361419.1361429.
- 3. Norberg, P.A., Horne, D.R. & Horne, D.A., 2007. The Privacy Paradox: Personal Information Disclosure Intentions versus Behaviors. *Journal of Consumer Affairs*, 41(1), pp.100-126.



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- The Marshall Symposium: Address Roger Needham, May 29, 1998, Rackham School of Graduate Studies, University of Michigan web.archive.org/web/20081201182254/http://www.si.umich.edu/marshall/doc s/p201.htm, accessed 2015-04-15.
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